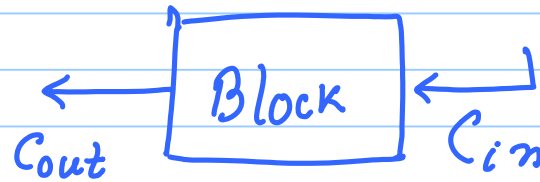


Aug-17



a_1
 b_1

	C_{in}	
	0	1
$C_{out,0}$	✓ ✓	✓
$C_{out,1}$	✓	✓ ✓

$$G \rightarrow a_1 \wedge b_1$$

$$P \rightarrow a_1 \oplus b_1$$

✓ → generate

✓ → absorb

✓ → propagate

$$C = \left[\begin{array}{|c|c|} \hline A & B \\ \hline \end{array} \right]^C$$

$\underbrace{\hspace{1.5cm}}_{P_A, G_A} \quad \underbrace{\hspace{1.5cm}}_{P_B, G_B}$

$$P_C = P_A \cdot P_B$$

$$G_C = G_A + G_B P_A$$

A		B
$a_1 \dots a_n$ + $b_1 \dots b_n$		$a_{n+1} \dots a_{n+m}$ $b_{n+1} \dots b_{n+m}$

	A		B
a	$a_{31} \dots a_{28}$		$a_{27} \dots a_0$
b	$b_{31} \dots b_{28}$		$b_{27} \dots b_0$

$$\begin{array}{r} a_1 a_0 \\ + b_1 b_0 \\ \hline \end{array}$$

$$P_0 = a_0 \oplus b_0$$

$$G_0 = a_0 \cdot b_0$$

$$P_1 = a_1 \oplus b_1$$

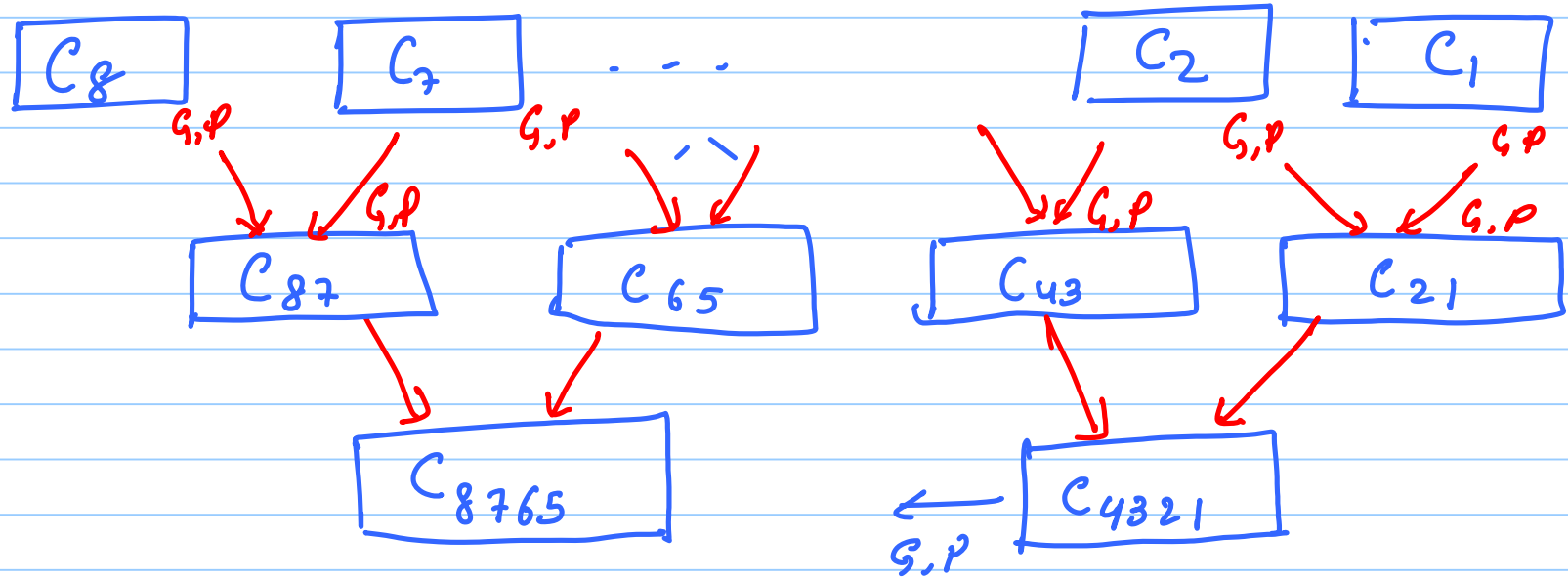
$$G_1 = a_1 \cdot b_1$$

$$P_{01} = P_0 \cdot P_1 = (a_0 \oplus b_0) (a_1 \oplus b_1)$$

$$G_{01} = G_1 + P_1 \cdot G_0 = a_1 \cdot b_1 + (a_1 \oplus b_1) \cdot a_0 \cdot b_0$$

Carry Lookahead Adder ($O(\log(n))$)

Round 1



(G, P of a block containing the first 16 bits)

Original
problem:

$$+ \begin{bmatrix} a_{31} & \dots & -a_{11} \\ b_{31} & \dots & -b_{11} \end{bmatrix} \begin{bmatrix} a_{15} & \dots & a_0 \\ b_{15} & \dots & b_0 \end{bmatrix}$$

$$C_{16} = G_{0-15}$$

$$+ \begin{bmatrix} a_{31} & \dots & a_{24} \\ b_{31} & \dots & b_{24} \end{bmatrix} \begin{bmatrix} a_{23} & \dots & a_{11} \\ b_{23} & \dots & b_{11} \end{bmatrix}$$

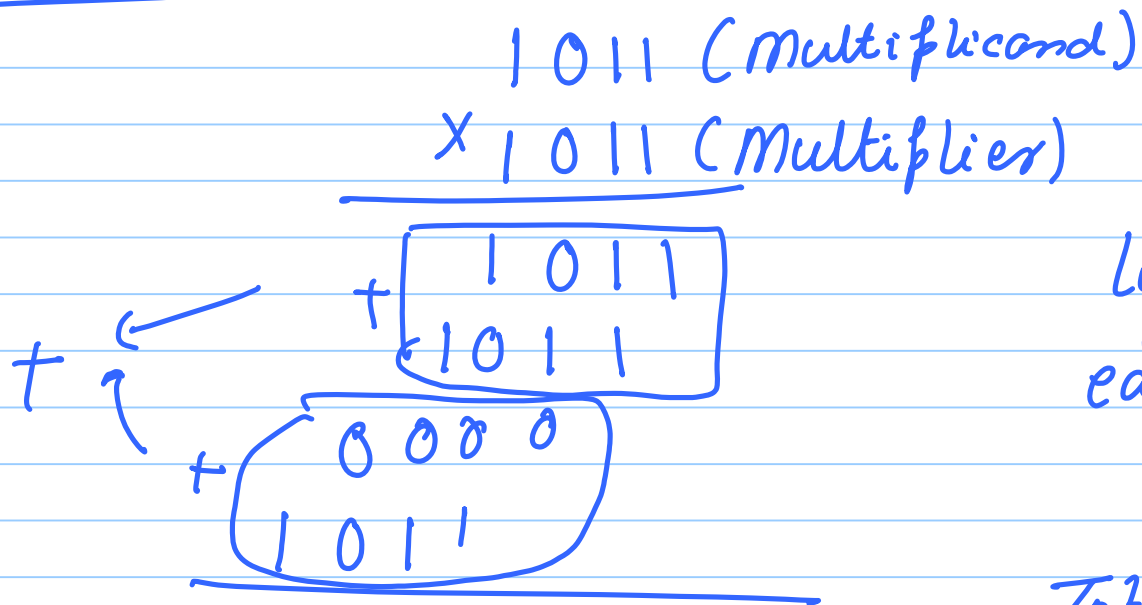
$$C_{24} = G_{16-23} + P_{16-23} C_{16}$$

ROUND - 2

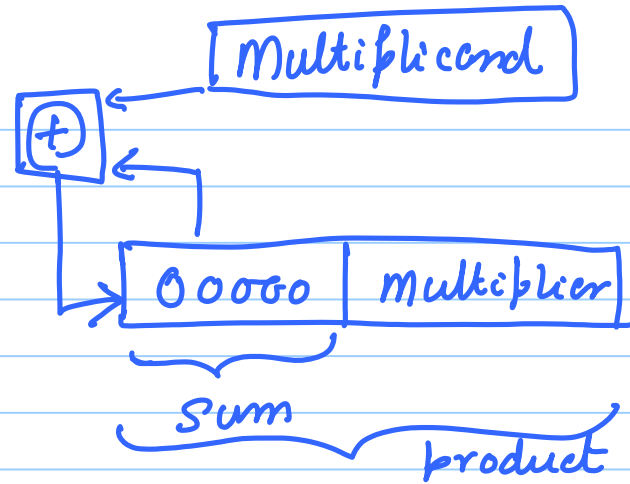
In Round 2:

You compute the C_{in} of each sub-block recursively until you reach the leaves. At that point use a ripple carry adder.

Multiplication



$\log(n)$ levels
each level takes
 $O(\log(n))$ time
Total: $O((\log(n))^2)$



repeat n times.

- 1) Take a look at LSB of multiplier.
- 2) If 0, go to step 4
- 3) If 1, $sum += multiplicand$
- 4) shift product 1 position right

Multiply two n bit numbers.

time: $O(n \log n)$ (Bad 😞)

Target: $(\log(n))^2$ -

Target for Wednesday: $O(\log(n))$