

Sept -26

Note Title

26-09-2011

Syllabus  $\rightarrow$  Uptil 21st

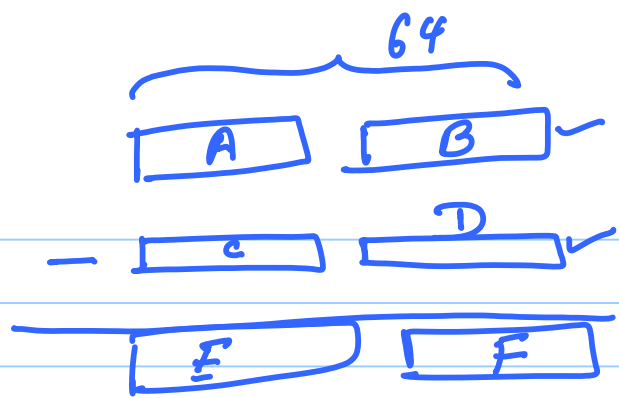
🚩 Class Notes

- Chapter 1-3  
(Book)

- Notes on Adder & FP

Focus:  $\Theta$  notation

Be very clear about  $\frac{O(\log(n))}{\text{what does it}}$



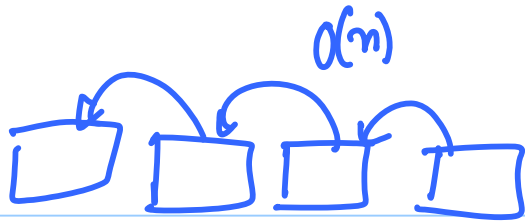
mean ?  
 Vacation  
4-9th

1

$\left\{ \begin{array}{l} \text{SUBS F, B, D} \\ \text{SBC E, A, C} \end{array} \right\} \quad (F = B - D)$

$\left\{ \begin{array}{l} F = B - D \\ E = A - C \\ \text{if } (D > B) \\ \quad E--; \end{array} \right.$

$D > B$   
 unsigned  
 comparison.



# Carry Lookahead Adder.

$$\begin{array}{r} a^{+cin} \\ + b \\ \hline \text{Cout Sum} \end{array}$$

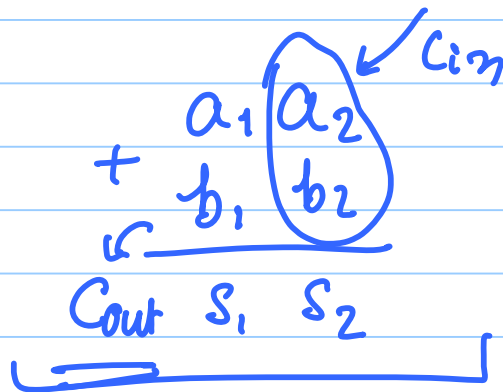
$$G = ab$$

$$\oplus \rightarrow \text{xOR}$$

$$P = (a \oplus b)$$

$$\text{Cout} = G + P \text{Cin}$$

(if Cin is 1  
Cout is 1)



$$G_2 = G_1 + P_1 G_2$$

$$P_2 = P_1 P_2$$

Invariant:

$$\text{Cout} = G + P \text{Cin}$$

$$\left\{ \begin{array}{cccc} a_k & a_{k-1} & \dots & a_1 \\ b_k & b_{k-1} & \dots & b_1 \end{array} \right\}$$

$$\underbrace{a_k \dots a_2}_{G_2, P_2} \quad \underbrace{a_{k-1} \dots a_1}_{G_1, P_1}$$

$$G = G_2 + P_2 G_1$$

$$P = P_1 P_2$$

$$\underbrace{a_k \dots a_2}_{G_{i,k}} \quad \underbrace{a_{k-1} \dots a_i}_{G_{i,i-1}}$$

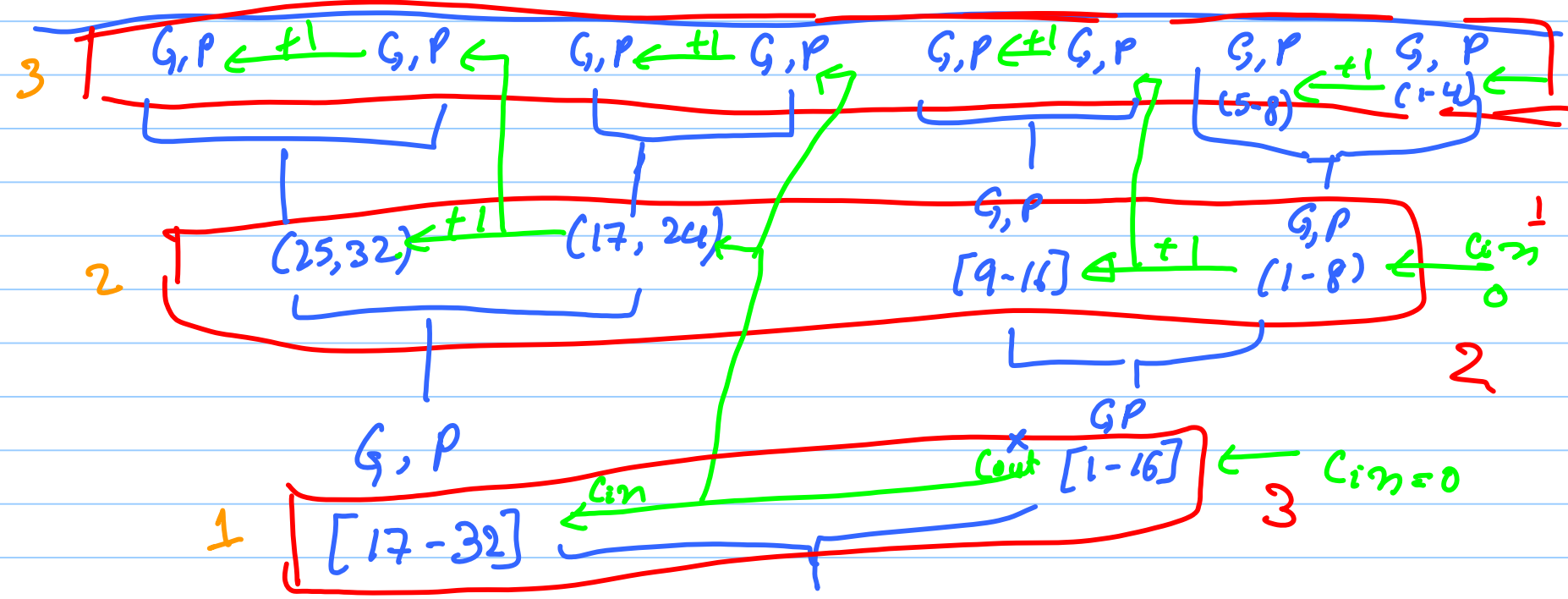
$$G_{i,k} = G_{i,k} + P_{i,k} \cdot G_{i,i-1}$$

$$P_{i,k} = P_{i,i-1} \cdot P_{i,k}$$

Invariant:

$$C_{out} = G + P C_{in}$$

4 bit blocks



$$[1-32] \\ \text{Cost} = G_{32} + P_{32} \cdot c_{in}^{(0)} = G_{(1,32)}$$

Stage I  
 $\log(n)$  steps: Generate & propagate signals (tree fashion)  $O(\log(n))$

Stage II: Compute the carry  
You have  $\rightarrow$  values of  $c_{in}$  for every block.  $\log(n)$

Stage III:

Do a ripple carry addition  
with the value of  $c_{in}$

$O(1)$

Total time :  $O(\log(n))$

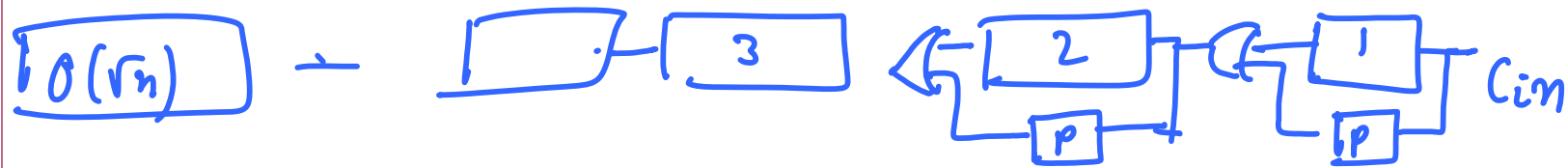
Question 6.

Only use Propagate

$g \equiv$  Cout of adding two blocks with

$c_{in} = 0$

CARRY-SKIP ADDER



stage I: Do the addition assuming  $C_{in} = 0$

$G + P \cdot C_{in} = C_{out}$   
 Do additions in parallel.

Output: All the generate functions  
 $O(\sqrt{n})$

Stage II: For each block sequentially

compute:  $C_{out} = G + P C_{in}$   
 $O(\sqrt{n})$



Output: The correct values of  $C_{in}$   
for every block

Stage II : Add again with correct value  
of  $C_{in}$ .

$O(\sqrt{n})$

Final Complexity :  $O(\sqrt{n})$

Next Class: Oct 10<sup>th</sup>

Final Exam: Nov. End